

On the periods of some graph transformations

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ABSTRACT

For any undirected graph with arbitrary integer values attached to the vertices, simultaneous updates are performed on these values, in which the value of a vertex is moved by 1 in the direction of the average of the values of the neighboring vertices. (A special rule applies when equality occurs.) It is shown these transformations always reach a cycle of length 1 or 2.

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1. Introduction

Let G be an undirected graph with vertices labelled $1, \dots, n$, and suppose that for each i , an integer $x_i(0)$ is initially assigned to vertex i . We perform a sequence of synchronous updates on these values. If $x_i(t)$ is the value of vertex i at time t , then:

$$x_i(t+1) = \begin{cases} x_i(t) - 1 & \text{if } \sum_{j \in J_i} x_j(t) < d_i x_i(t), \\ x_i(t) & \text{if } \sum_{j \in J_i} x_j(t) \geq d_i x_i(t) \text{ and} \\ & \text{for some } j \in J_i, x_j(t) \neq x_i(t), \\ x_i(t) + 1 & \text{if } \sum_{j \in J_i} x_j(t) < d_i x_i(t), \\ & \text{for all } j \in J_i, x_j(t) = x_i(t) \end{cases}$$

where

$$J = \{j: \text{vertex } j \text{ is connected to vertex } i\},$$

$$d = |J_i| = \text{degree of vertex } i.$$

Less formally, the value $x_i(t)$ assigned to vertex i moves by 1 in the direction of the average of the values assigned to the neighbors of vertex i , but a special rule applies when $x_i(t)$ equals this average. Since $\max_i x_i(t)$ does not increase and $\min_i x_i(t)$ does not decrease as t varies, the iteration described above eventually reaches a cycle, so that

for some minimal $p > 1$, $x_i(t+p) = x_i(t)$ for all i and all $t > t_0$. Our main result is that the length of the cycle is 1 or 2.

Theorem: For any undirected graph G and any initial assignment of integers $x_1(0), \dots, x_n(0)$ to the vertices of G , there is a t_0 such that the above iteration satisfies $x_i(t+2) = x_i(t)$ for all i and all $t > t_0$.

The problem of determining the cycle length of the above iteration arose in the work of D. Ghiglia and G. Mastin [1]. They considered such iterations for the cases of G being (a) a simple path (i.e., vertex i being connected to vertex j if and only if $|i-j| = 1$) and (b) a k by m rectangular grid of lattice points, with edges between points that are horizontal or vertical neighbors. The rules described above were constructed as part of an algorithm for “phase unwrapping”; i.e., determining the argument of a complex function given the principal value of the argument, so on to eliminate the discontinuation by integer multiples of 2π .

The “phase unwrapping” origin of the transformation accounts for the irregularity in the rules prescribing that if the average of the values of a sites neighbors equals the value at that site, but not all the neighbors are equal to the site, then the value of the site should be incremented by 1. As it turns out, even if this condition is relaxed, the length of the cycle is still at most 2. The proof of this is similar to that of our theorem.

Ghiglia and Mastin found by extensive simulations that iterations of the transformation always led to cycles of length 1 or 2. They conjectured that this is always the case, and their “phase unwrapping” algorithm is based on the assumption that this conjecture is true. Our theorem, which proves this conjecture, guarantees that the

Ghighlia-Mastin algorithm will always terminate.

E. Brickell and M. Purtil were the first to consider the general transformation as we defined it above. When all the $x_i(0)$ are 0 or 1, they showed (unpublished) assigned a value of by a very elegant combinatorial argument that the cycle length is at most two.

At any time t , divide the vertices of G into four classes as follows:

$$c_1 = \{i: x_i(t) = 0 \text{ and } x_j(t) = 0 \text{ for all } j \in J_i\},$$

$$c_2 = \{i: x_i(t) = 1 \text{ and } x_j(t) = 1 \text{ for all } j \in J_i\},$$

$$c_3 = \{i: x_i(t) = 0 \text{ and there exists } j \in J_i \text{ with } x_j(t) = 1\},$$

$$c_4 = \{i: x_i(t) = 1 \text{ and there exists } j \in J_i \text{ with } x_j(t) = 1\}.$$

Any site in c_1 at time t will be in c_1 or c_3 at time $t+1$ since the value will remain 0, but we cannot predict what will happen to its neighbors. Similarly, any site which falls in c_2 at time t will be in c_2 or c_4 at time $t+1$. Anything in c_3 will move to c_4 at time $t+1$, and all members of c_4 will move to c_3 . Therefore eventually all elements will either stay in c_1 or in c_2 or will continue switching between c_3 and c_4 , and so the cycle length will be 1 or 2.

Where the $x_i(0)$ are not all 0 or 1 (or x and $x+1$, more generally), the iteration is much more complicated and no simple combinatorial argument has been found to prove the theorem. For example, even when G is a simple path, differences between values of adjacent vertices can be arbitrarily large on a cycle (as large as a constant times n for a path of length n). This can be seen by generalizing the construction $(x_1(0), \dots, x_H(0)) = (0, 1, 1, 4, 6, 11, 15, 22, 25, 28, 27)$.

The proof we will give for the theorem is based on a modification of the proof used by Goles-Chacc, Fogelman-Soulic, and Pellegrin [2] to prove that cycle lengths are at most two in certain threshold networks. Their Theorem implies the Brickell-Purtill results, but does not seem to cover the general case of our iteration. However, their concept of decreasing energy is a key ingredient in our proof. Another case where iterations on graphs produce cycles of length at most 2 occurs in the work of Poljanc and Sora [3], but their model and method of proof are quite different from ours.

2. Proof of theorem

It clearly suffices to prove the theorem when G is connected, and so we will assume this from now on.

Lemma: If the period of the cycle is not 1, then for all large t and for all i ,

$$x_i^t(t) \neq x_i(t+1) .$$

Proof of lemma: Suppose there exists i' , t such that $x_{i'}(t) = x_{i'}(t+1)$ and that the t -th iteration is in the cycle. We know there exists j' such that $x_{j'}(t) \neq x_{j'}(t+1)$ since the period of the cycle is not 1. Hence we can find i and j that are connected such that:

$$\text{but } x_i(t) = x_i(t+1) ,$$

$$x_j(t) \neq x_j(t+1) ,$$

$$x_j(t+1) = x_j(t) \pm 1 .$$

$$\text{So } x_i(t) - x_j(t) \equiv 0 \pmod{2}$$

and

$$x_i(t+1) - x_j(t+1) \equiv 1 \pmod{2} .$$

But if $x_i(t+k) - x_j(t+k) \equiv 1 \pmod{2}$, then $x_i(t+k) \neq x_j(t+k)$, so $x_i(t+k+1) = x_i(t+k) \pm 1$ and $x_j(t+k+1) = x_j(t+k) \pm 1$, so $x_i(t+k+1) - x_j(t+k+1) \equiv 1 \pmod{2}$.

Since this is true for all k , there does not exist any $k' > 0$ such that $x_i(t+k') = x_j(t+k')$ which means that $x_i(t)$ cannot be in the cycle and we have reached a contradiction, which proves the lemma. ■

Proof of Theorem: Using the idea of a decreasing “energy” function utilized by Goles-Chacc et. al. [2], we define:

$$E(t) = - \sum_{i,j=1}^n a_{ij} x_i(t) x_j(t-1),$$

where

$$a_{ij} = \begin{cases} 1 & \text{if } i \neq j \text{ but } j \in J_i; \\ -d_i & \text{if } i = j; \\ 0 & \text{if } i \neq j \text{ and } j \notin J_i. \end{cases}$$

Note that E is bounded below since the maximal element at any stage never increases with time.

We now consider the change in energy during iterations of the transformation:

$$\begin{aligned} \Delta E(t) &= E(t+1) - E(t) = - \sum_{i,j=1}^n a_{ij} x_i(t+1) x_j(t) - \sum_{i,j=1}^n a_{ij} x_i(t) x_j(t-1) \\ &= - \sum_{i=1}^n [(x_i(t+1) - x_i(t-1)) \sum_{j=1}^n a_{ij} x_j(t)], \quad (*) \end{aligned}$$

since $a_{ij} = a_{ji}$ for all i, j . For each i , if $\sum_{j=1}^n a_{ij} x_j(t) < 0$, then

$$d_i x_i(t) > \sum_{j \in J_i} x_j(t),$$

so

$$x_i(t+1) < x_i(t),$$

$$x_i(t+1) - x_i(t-1) \leq 0,$$

$$\text{and } -(x_i(t+1) - x_i(t-1)) \sum_{j=1}^n a_{ij} x_j(t) \leq 0.$$

If $\sum_{j=1}^n a_{ij} x_j(t) \geq 0$, then:

$$d_i x_i(t) \leq \sum_{j \in J_i} x_j(t),$$

$$x_i(t+1) > x_i(t),$$

$$x_i(t+1) - x_i(t-1) \geq 0,$$

$$\text{and } -(x_i(t+1) - x_i(t-1)) \sum_{j=1}^n a_{ij} x_j(t) \leq 0.$$

Thus in both cases $\Delta E(t) \leq 0$ and each term in the sum on i on the right side of (*) is ≤ 0 . Since E is bounded below, we must have $\Delta E(t) = 0$ for all large $t \geq t_0$, and moreover, for all such $t \geq t_0$ we have for each i ,

$$(x_i(t+1) - x_i(t-1)) \sum_{j=1}^n a_{ij} x_j(t) = 0.$$

Now suppose there exists an i such that $x_i(t+1) \neq x_i(t-1)$ for some $t \geq t_0$. If $x_i(t-1) > x_i(t+1)$ and

$$\sum_{j \in J_i} x_j(t) > d_i x_i(t),$$

then

$$x_i(t+1) > x_i(t),$$

which is a contradiction. If

$$x_i(t+1) = x_i(t-1)$$

then there exists a k such that $x_i(t+k+1) > x_i(t+k-1)$, so

$$\sum_{j \in J_i} x_j(t+k) > d_i x_i(t+k),$$
 which is also a contradiction.

Therefore $x_i(t+1) = x_i(t-1)$ whenever $x_i(t-1)$ is in the cycle, so the length of the cycle is at most 2. ■

Note: It is not true that $\Delta E(t) < 0$ for t not in the cycle. For example, when G consists of a simple path of length 5, with $(x_1(0), \dots, x_5(0)) = (0, 2, 2, 3, 5)$, then $E(1) = E(2) = -1$, but the cycle starts only at $t = 2$.

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